

5.7 Searching

Name: _____ Class: _____ Date: _____

Total: 24 marks**KEY WORDS** linear search • binary search

Objective

Run both searches by hand and respect the sorted precondition.

1 Key knowledge and worked examples

■ The two searches

Linear walks, $O(n)$, no preconditions; binary halves a sorted range – 16 slots resolve in at most $\lceil \log_2 17 \rceil = 5$ probes ✓. On the course array, finding 43 probes 7, 11, 9, 10 ✓.

■ Worked example: the miss

Searching 64 probes 7(31), 11(47), 13(55), 14(59), 15(63) – five probes, then the empty range reports absent; linear walks all 16 ✓.

2 Practice

2.1 A sorted array of 32 slots: give the worst-case binary probe count. [3]

2.2 State the update that makes binary search loop forever. [3]

2.3 Give the safe midpoint formula and the failure it prevents in fixed-width languages. [3]

3 Exam-style questions

3.1 The sorted array is $[2, 6, 10, 14, 18, 22, 26, 30]$ (indices 0–7).

Search 22: give every probe index, the value read, and the range kept after each probe.
Search 5 (absent): give the probe trail and the final empty range. Give both worst-case bounds (8 slots) and the exact linear cost of each query. [8]

3.2 An unsorted array of 10^6 values must answer 10^4 queries.

Price all-linear: total comparisons. Price sort-then-binary: the $\Theta(n \log_2 n)$ sort cost (value at $n = 10^6$), plus the 10^4 queries at at most $\lceil \log_2(10^6 + 1) \rceil$ probes each (value of the ceiling). State the crossover condition on the number of queries that justifies sorting first. [7]

Solutions

2.1

- $\lceil \log_2 33 \rceil = 6$ probes ✓.

2.2

- `lo = mid` (instead of `mid + 1`) with `hi` unchanged loops forever ✓.

2.3

- `lo + (hi - lo) // 2` ✓.
- Prevents `lo + hi` overflowing a fixed-width integer ✓.

3.1

- 22: mid 3 reads $14 < 22 \rightarrow lo=4$; mid 5 reads 22 – found: probes 3, 5 ✓.
- 5: mid 3 reads $14 > 5 \rightarrow hi=2$; mid 1 reads $6 > 5 \rightarrow hi=0$; mid 0 reads $2 < 5 \rightarrow lo=1 > hi=0$ – empty, absent ✓.
- Worst-case bounds (8 slots): $\lceil \log_2 9 \rceil = 4$ probes ✓.
- Linear: 22 sits at index 5 – 6 values; the miss 5 walks all 8 ✓.

3.2

- All-linear: $10^4 \times 10^6 = 10^{10}$ comparisons ✓.
- Sort: $\Theta(n \log_2 n) = 10^6 \times 19.93 \approx 1.99 \times 10^7$ comparisons ✓.
- Queries: $\lceil \log_2(10^6 + 1) \rceil = 20$ probes each – 2×10^5 total ✓.
- Sort-then-binary total $\approx 2 \times 10^7$, about 500 times less than 10^{10} ✓.
- Crossover: sorting wins once query count q satisfies $q \cdot n > n \log_2 n + q \log_2 n$ – roughly $q > \log_2 n \approx 20$ ✓.