

5.1 Asymptotic analysis

Name: _____ Class: _____ Date: _____

Total: 24 marks

KEY WORDS Big-O O • Big-Θ Θ

Objective

Classify algorithms by growth, and prove the class by doubling.

1 Key knowledge and worked examples

■ The classes

Drop constants and lower-order terms: $3n^2 + 5n + 7$ is $\Theta(n^2)$ ✓. At $n = 1000$: $n \log_2 n \approx 9\,966$ against $n^2 = 10^6$ – a factor of about 100 ✓.

■ Worked example: doubling reads the exponent

Nested loops at $n = 100$ run their body 10 000 times; at $n = 200$, 40 000 ✓. The ratio 4 is the exponent: the same doubling turns a linear loop's 100 into 200 ✓.

2 Practice

2.1 State the class of $7n \log_2 n + 2n + 40$, naming the rule you applied. [3]

2.2 A halving loop starts at $n = 1024$. State how many iterations to reach 1. [3]

2.3 Give the $\log_2$ values for $n = 10, 100, 1000$, rounded to integers. [3]

3 Exam-style questions

3.1 Two programs run on the same machine. Program A performs $100n$ operations; program B performs $2n^2$ operations.

State the size n at which B first overtakes A (solve $2n^2 = 100n$). Classify each program, justifying by the dominant term. State which program you run at $n = 10$ and which at $n = 1000$, with the operation counts at both sizes for each program. [8]

3.2 A loop runs `while i > 1: i = i // 2` and inside it a second loop runs n times.

State what each loop contributes and multiply them into one class. Compare its cost at $n = 1024$ with a pure n^2 loop at the same n , computing both values. Explain in one line why halving produces a logarithm. [7]

Solutions

2.1

- $\Theta(n \log n)$ ✓.
- Dominant term wins; constants and $2n$ dropped ✓.

2.2

- 10 iterations: $2^{10} = 1024$ ✓.

2.3

- 3, 7, 10 ✓.

3.1

- $2n^2 = 100n$ gives $n = 50$: B first overtakes A there ✓.
- A is $\Theta(n)$ (linear); B is $\Theta(n^2)$ – the dominant terms decide ✓.
- At $n = 10$: A 1 000, B 200 – run B ✓.
- At $n = 1000$: A 100 000, B 2 000 000 – run A ✓.

3.2

- The halving loop contributes $\log_2 n$ iterations; the inner loop n each time ✓.
- Product: $\Theta(n \log n)$ ✓.
- At $n = 1024$: $n \log_2 n = 1024 \times 10 = 10\,240$ against $n^2 = 1024^2 = 1\,048\,576$ – about 102 times more ✓.
- Each iteration halves the input, so the count of steps to reach 1 is the number of halvings – the logarithm ✓.