

Data Collections

AP Computer Science A ■ Unit 4

AP Computer Science A

4	8	2	9
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Data Collections

What we will cover

- 1 Ethics of collecting data
- 2 Arrays
- 3 ArrayList
- 4 2D arrays (grids)
- 5 Searching & sorting
- 6 Recursion

4	8	2	9	5
0	1	2	3	4

1 Ethics

Data Collections

└ Ethics

The ethics of collecting data

Programs that collect and process personal data carry real responsibility.

Collect only what you need, protect **privacy** 隐私, be honest about how data is used, and guard against **bias** 偏差 in what you gather. Legal is not always ethical.



Data centre

2 Arrays

Data Collections

└ Arrays

Arrays

4	8	2	9	5	7
0	1	2	3	4	5

An **array** 数组 stores a fixed number of same-type values, reached by **index** 索引 (from **0**). Visit every element with a loop:

```
for (int i = 0; i < a.length; i++) ...
```

Standard algorithms: find the **max/min**, **sum/average**, or count matches by traversing once.



Filing cabinet

3 ArrayList

Data Collections
└ ArrayList

The ArrayList toolbox

An **ArrayList** 动态数组 grows and shrinks, unlike a fixed array. It stores objects, so numbers are **wrapped** 包装 (Integer, Double).

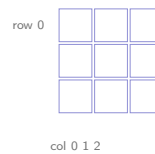
add(x) append	get(i) read	set(i,x) replace	remove(i) delete
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Traverse it with a for or an enhanced for-each loop: `for (int x : list).`

4 2D arrays

Data Collections
└ 2D arrays

2D arrays (grids)



A **2D array** 二维数组 is a grid of rows & columns: `grid[r][c]`.

Walk it with **nested loops** – outer over rows, inner over columns:

```
for r ...  
  for c ...  
    grid[r][c]
```

5 Search & sort

Data Collections
└ Search & sort

Searching & sorting

Linear search 线性查找
checks every element; **binary search** 二分查找
halves a **sorted** array

Selection sort 选择排序 & **insertion sort** 插入排序 order a list;
merge sort 归并排序 is faster

Binary search needs **sorted** data and does $\sim \log_2 n$ comparisons, far fewer than linear search's n .

Data Collections └ Search & sort Recursion A recursive 递归 method calls itself on a smaller input, stopping at a base case 基本情形. Factorial $\text{factorial}(n) = n * \text{factorial}(n-1)$, with $\text{factorial}(0) = 1$. Each call shrinks n until the base case, then the results multiply back up. Merge sort 归并排序 uses the same divide-and-conquer idea.	Data Collections └ Search & sort Key points Arrays are fixed; ArrayLists grow & hold objects A 2D array is a grid, walked with nested loops Binary search & merge sort beat the simple versions Recursion solves a problem via a smaller copy
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Data Collections └ Search & sort Worked example – tracing recursion Trace factorial(4) Each call defers to a smaller one: $4 \cdot f(3) = 4 \cdot 3 \cdot f(2) = 4 \cdot 3 \cdot 2 \cdot f(1)$ $\text{factorial}(1)$ hits the base case and returns 1, so the calls unwind inward: $2 \cdot 1 = 2$, then $3 \cdot 2 = 6$, then $4 \cdot 6 = 24$. Writing each call above its returned value is the reliable way to trace recursion – and the base case is what stops it.	Data Collections └ Search & sort Check yourself 1 factorial(4) returns what? 2 What happens to recursion with no base case? 3 Binary search needs the array to be ...? 4 Array or ArrayList: which resizes itself? Answers 1. 24 2. it never stops – a StackOverflowError 3. sorted 4. ArrayList
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