

Data Collections

AP Computer Science A ■ Unit 4

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4	8	2	9
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What we will cover

1 Ethics of collecting data

2 Arrays

3 ArrayList

4 2D arrays (grids)

5 Searching & sorting

6 Recursion

4	8	2	9	5
0	1	2	3	4

The ethics of collecting data

Programs that collect and process personal data carry real responsibility.

Collect only what you need, protect **privacy** 隐私, be honest about how data is used, and guard against **bias** 偏差 in what you gather. Legal is not always ethical.



data centre

Arrays

4	8	2	9	5	7
0	1	2	3	4	5

An **array** 数组 stores a fixed number of same-type values, reached by **index** 索引 (from **0**). Visit every element with a loop:

```
for (int i = 0; i < a.length; i++) ...
```

Standard algorithms: find the **max/min**, **sum/average**, or count matches by traversing once.



filing cabinet

Arrays

index	1	2	3	4	5
scores	90	75	60	88	50

scores[2] is 75

The ArrayList toolbox

An **ArrayList** 动态数组 grows and shrinks, unlike a fixed array. It stores objects, so numbers are **wrapped** 包装 (Integer, Double).

add(x)
append

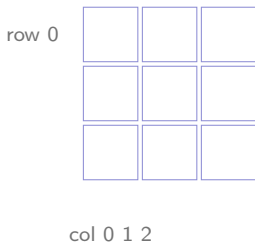
get(i)
read

set(i,x)
replace

remove(i)
delete

Traverse it with a for or an enhanced for-each loop: `for (int x : list).`

2D arrays (grids)



A **2D array** 二维数组 is a grid of rows & columns:
`grid[r][c]`.

Walk it with **nested loops** – outer over rows, inner over columns:

```
for r ...  
    for c ...  
        grid[r][c]
```

2D arrays (grids)

		column		
		1	2	3
row	1			
	2			
	3			

← grid[2,3]

Searching & sorting

Linear search 线性查找
checks every element; **binary search 二分查找**
halves a **sorted** array

Selection 选择排序 & insertion sort 插入排序 order a list;
merge sort 归并排序 is faster

Binary search needs **sorted** data and does $\sim \log_2 n$ comparisons, far fewer than linear search's n .

Searching & sorting



Recursion

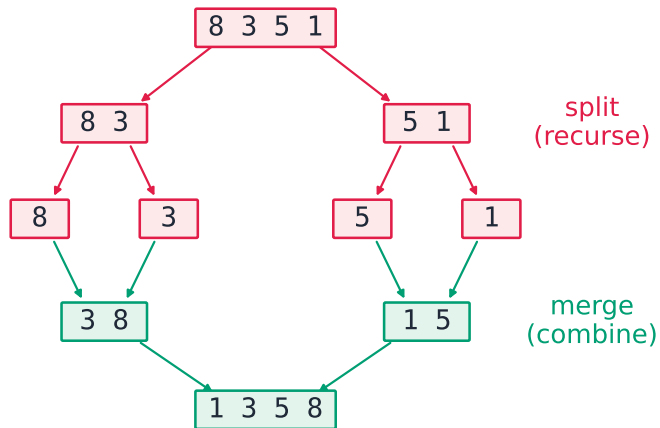
A **recursive** 递归 method calls itself on a smaller input, stopping at a **base case** 基本情形.

Factorial

$\text{factorial}(n) = n * \text{factorial}(n-1)$, with $\text{factorial}(0) = 1$.

Each call shrinks n until the base case, then the results multiply back up. **Merge sort** 归并排序 uses the same divide-and-conquer idea.

Recursion



Key points

Arrays are fixed; **ArrayLists** grow & hold objects

A **2D array** is a grid,
walked with nested loops

Binary search & merge sort
beat the simple versions

Recursion solves a problem
via a smaller copy

Worked example – tracing recursion

Trace factorial(4)

Each call **defers** to a smaller one:

$$4 \cdot f(3) = 4 \cdot 3 \cdot f(2) = 4 \cdot 3 \cdot 2 \cdot f(1)$$

factorial(1) hits the **base case** and returns 1, so the calls **unwind inward**: $2 \cdot 1 = 2$, then $3 \cdot 2 = 6$, then $4 \cdot 6 = 24$.

Writing each call **above** its returned value is the reliable way to trace recursion – and the base case is what stops it.

Check yourself

- 1 `factorial(4)` returns what?
- 2 What happens to recursion with no base case?
- 3 Binary search needs the array to be ...?
- 4 Array or ArrayList: which resizes itself?

Answers 1. 24 2. it never stops – a `StackOverflowError` 3. sorted 4. ArrayList