

19. Computational thinking and Problem-solving Starter Quiz · 课前小测

A-Level Computer Science

Name: _____ Score: _____

1. A linear search:

- ☐ checks each element in turn and works on any list
- ☐ requires the data to be sorted
- ☐ always finds the target in one step
- ☐ only works on numbers

2. Bubble sort works by:

- ☐ repeatedly swapping adjacent out-of-order pairs
- ☐ inserting each element into a sorted prefix
- ☐ halving the list each time
- ☐ building a binary tree

3. Which Big-O describes binary search?

- ☐ $O(\log n)$
- ☐ $O(n)$
- ☐ $O(n^2)$
- ☐ $O(n \log n)$

4. Every recursive algorithm must have a base case because:

- ☐ without it the recursion never stops
- ☐ it makes the code shorter
- ☐ it speeds up the program
- ☐ the compiler requires a comment

5. The worst-case time complexity of binary search is:

- ☐ $O(\log n)$
- ☐ $O(n)$
- ☐ $O(n^2)$
- ☐ $O(1)$

6. What does Factorial(4) return?

7. About how many comparisons does a binary search need for one million sorted items?

8. Insertion sort runs close to $O(n)$ on small or nearly-sorted arrays, because few elements need to be shifted.

☐ True ☐ False

9. A stack (LIFO) naturally drives a depth-first traversal, while a queue (FIFO) drives a breadth-first traversal.

☐ True ☐ False

10. Recursion is a natural fit for self-similar problems (trees, divide-and-conquer), but each call adds a stack frame — so without a base case it overflows the stack.

☐ True ☐ False

A-Level Computer Science

1. **checks each element in turn and works on any list**

Linear search needs no preparation and works on any list, at worst $O(n)$.

2. **repeatedly swapping adjacent out-of-order pairs**

Each pass swaps adjacent pairs, "bubbling" the largest element to the end.

3. **$O(\log n)$**

Halving the range each step is logarithmic — $O(\log n)$.

4. **without it the recursion never stops**

The base case is the condition that ends the chain of calls; without it the recursion runs forever.

5. **$O(\log n)$**

Halving the range each step gives a logarithmic number of comparisons.

6. **24**

$$4 \times 3 \times 2 \times 1 = 24.$$

7. **20**

$$\log_2(1\,000\,000) \approx 20 \text{ — about 20 comparisons.}$$

8. **True**

On nearly-sorted data each new item is already almost in place — which is why insertion sort beats fancier sorts on small inputs.

9. **True**

The ADT you choose decides the search order — stack goes deep first, queue explores level by level.

10. **True**

A simple counting loop is cleaner for plain iteration; recursion shines when the problem contains smaller copies of itself.